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Introduction to Cryptography

EECE 412

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In the News: Hacker exposes embassies' e-mail

A Swedish security expert released last week the addresses and passwords for 100 e-mail accounts, claiming that he has uncovered a flaw that exposes more than a thousand sensitive e-mail accounts at government agencies, such as embassies, and corporations.

The e-mail account information appeared on the DerangedSecurity blog, run by Swedish hacker Dan Egerstad, and listed the e-mail server IP addresses, e-mail addresses, and passwords for accounts at numerous embassies, including the Russian, Indian, and Iranian embassies in various countries. Other accounts belonged to government officials and civil-rights workers. While Egerstad released the information for 100 accounts, he told Wired News that he had collected more than 1,000.

"Here is everything you need to read classified email and f**k up some serious international business," wrote Egerstad on his blog. "Hopefully this will put light on the security problems that are never talked about and get at least this fixed with a speed that you never seen your government work before."

Outing the poor security of government agencies has its risks. In 2006, the FBI raided the home of a security researcher that pointed out the insecurities in boarding pass checks, and created a Web site to allow people to print out their own passes. In 2003, authorities arrested Brett E. O'Keefe, president of California start-up ForensicTec, after he demonstrated the insecurities in several U.S. military networks by hacking into them. Two years later, O'Keefe was sentenced to 60 days in a work release program.

In the latest incident, Egerstad decided not to notify each organization because he did not believe that they would listen. He also admitted to viewing thousands of classified e-mails.

Source: securityfocus.com
2007-09-04



Session Outline

- Historical background
 - Classic ciphers
 - One-time pad
 - One-way functions
- The Random Oracle model
 - Random functions: Hash functions
 - Random generators: stream ciphers
 - Random Permutations: block ciphers
 - Public key encryption and trapdoor one-way permutations
 - Digital signatures



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Crypto

- **Cryptology** — The art and science of making and breaking “secret codes”
- **Cryptography** — making “secret codes”
- **Cryptanalysis** — breaking “secret codes”
- **Crypto** — all of the above (and more)



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How to Speak Crypto

- A *cipher* or *cryptosystem* is used to *encrypt* the *plaintext*
- The result of encryption is *ciphertext*
- We *decrypt* ciphertext to recover plaintext
- A *key* is used to configure a cryptosystem
- A *symmetric key* cryptosystem uses the same key to encrypt as to decrypt
- A *public key* cryptosystem uses a *public key* to encrypt and a *private key* to decrypt



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Crypto

- Basis assumption
 - The system is completely known to the attacker
 - Only the key is secret
- Also known as **Kerckhoffs Principle**
 - Crypto algorithms are not secret



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Kerckhoff's Principle

“The security of a cryptosystem must not depend on keeping secret the crypto-algorithm. The security depends only on keeping secret the key”

Auguste Kerckhoff von Nieuwenhof

Dutch linguist

1883



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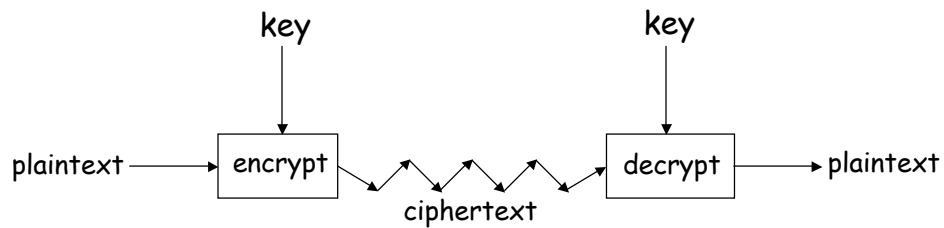
Crypto

- Basis assumption
 - The system is completely known to the attacker
 - Only the key is secret
- Also known as **Kerckhoffs Principle**
 - Crypto algorithms are not secret
- Why do we make this assumption?
 - Experience has shown that secret algorithms are weak when exposed
 - Secret algorithms never remain secret
 - Better to find weaknesses beforehand



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Crypto as Black Box



A generic use of crypto



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Historical Background

To read:

5.1-5.2 Anderson's book

Chapter 2 (except 2.3.6 & 2.3.8)

Stamp's book

Letter Indices in English Alphabet

A	B	C	D	E	F	G	H	I	J	K	L	M
0	1	2	3	4	5	6	7	8	9	10	11	12
N	O	P	Q	R	S	T	U	V	W	X	Y	z
13	14	15	16	17	18	19	20	21	22	23	24	25

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Caesar Cipher

- Plaintext is HELLO WORLD
- Change each letter to the third letter following it (X goes to A, Y to B, Z to C)
 - Key is 3, usually written as letter 'D'
 - $C = P + K \text{ mod } 26$

• Ciphertext: KHOOR ZRUOG

Plain HELLOWORLD

Key DDDDDDDDDD

Cipher KHOORZRUOG

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Monoalphabetic Substitution Cipher

Invented by Arabs in 8th or 9th centuries

A	B	C	D	E	F	G	H	I	J	K	L	M	N	..	Z
F	T	W	S	G	M	P	A	Z	C	L	V	O	D	..	B

Plain HELLOWORLD

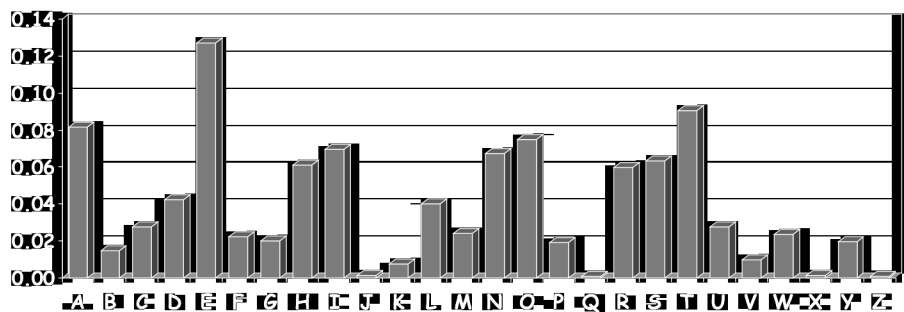
Key

Cipher AGVVYEYZVS



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Frequency Analysis



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Polyalphabetic Vigenère Cipher

proposed by Blaise de Vigenere from the court of Henry III of France in the sixteenth century

Like Cæsar cipher, but use a phrase

▪ Example

- Message: TO BE OR NOT TO BE THAT IS THE QUESTION
- Key: RELATIONS
- Encipher using Cæsar cipher for each letter:

Plain TO BE OR NOT TO BE TH AT IS THE QUESTION
 Key RE LA T I ONS RE LA T I ON SR ELA T I ONSREL
 Cipher KS ME HZ BBL KS ME MPOG AJ XSE J CSFLZSY



Cryptanalysis of Vigenère Cipher

Factoring of distances

- KSMEHZBBLKSMEMPOGAJXSEJCSFLZSY
- 012345678012345678012345678012



Cryptanalysis: Terminology

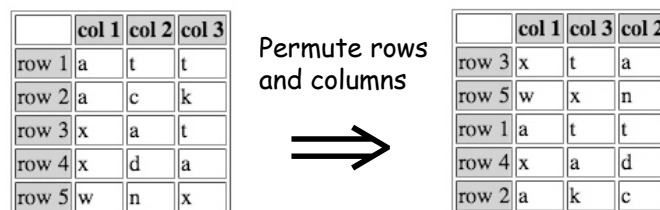
- Cryptosystem is **secure** if best know attack is to try all keys
- Cryptosystem is **insecure** if any shortcut attack is known
- By this definition, an insecure system might be harder to break than a secure system!

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Double Transposition

- Plaintext: attackatxdawn



- Ciphertext: xtawxnattxadakc
- Key: matrix size and permutations (3,5,1,4,2) and (1,3,2)

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One-Time Pad

A Vigenère cipher with a random key at least as long as the message

- Provably unbreakable
- Why?

Plain text	D O I T	D O N T
Key	A J I Y	A J D Y
Cipher text	D X Q R	D X Q R

- Warning: keys *must* be random, or you can attack the cipher by trying to regenerate the key



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Little Bit of History

90 years ago,
January 19, 1917 ...

Codebook

- Literally, a book filled with “codewords”
- Zimmerman Telegram encrypted via codebook

Februar	13605
fest	13732
finanzielle	13850
folgender	13918
Frieden	17142
Friedenschluss	17149
:	:

- Modern block ciphers are codebooks!

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Zimmerman Telegram

- One of most famous codebook ciphers ever
- Led to US entry in WWI
- Ciphertext shown here...

WESTERN UNION TELEGRAM

GERMAN LEGATION
MEXICO CITY

via Galveston
JAN 19 1917

130	13042	13401	8501	115	3528	415	17214	0491	11310
18147	18222	21560	10247	11518	23077	13605	3494	14956	
98092	5905	11311	10392	10371	0302	21290	5101	39095	
23571	17504	11099	18276	18101	0317	0228	17694	4473	
22224	22200	19452	21589	07893	5569	15918	8958	12137	
1333	4725	4458	5905	17105	13851	4458	17149	14471	0706
13850	12224	0929	14991	7382	15857	07895	14218	36477	
5870	17553	07892	5870	5454	16102	15217	22801	17138	
21001	17388	7446	23038	18222	0719	14331	15021	23845	
3156	23552	22096	21804	4797	9497	22461	20855	4377	
23610	18140	22260	5905	13347	20420	39689	13732	20607	
0929	5275	18507	52962	1340	22049	13339	11265	22295	
10439	14814	4178	0592	8784	7632	7357	0928	52282	11287
21100	21272	9346	9559	22474	15874	18502	18500	15857	
2188	5376	7381	98092	16127	13486	9350	9220	76036	14219
5144	2831	17520	11347	17142	11264	7867	7782	15099	9110
10482	97556	3569	3670						

BEHNSTOFF.

Charge German Embassy.

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Zimmerman Telegram Decrypted

- British had recovered partial codebook
- Able to fill in missing parts

TELEGRAM RECEIVED.
 FROM 2nd from London # 5747.
 "We intend to begin on the first of February unrestricted submarine warfare. We shall endeavor in spite of this to keep the United States of America neutral. In the event of this not succeeding, we make Mexico a proposal of alliance on the following basis: make war together, make peace together, generous financial support and an understanding on our part that Mexico is to reconquer the lost territory in Texas, New Mexico, and Arizona. The settlement in detail is left to you. You will inform the President of the above most secretly as soon as the outbreak of war with the United States of America is certain and add the suggestion that he should, on his own initiative, ~~invite~~ ^{invite} Japan to immediate adherence and at the same time mediate between Japan and ourselves. Please call the President's attention to the fact that the ruthless employment of our submarines now offers the prospect of compelling England in a few months to make peace." Signed, ZIMMERMAN.

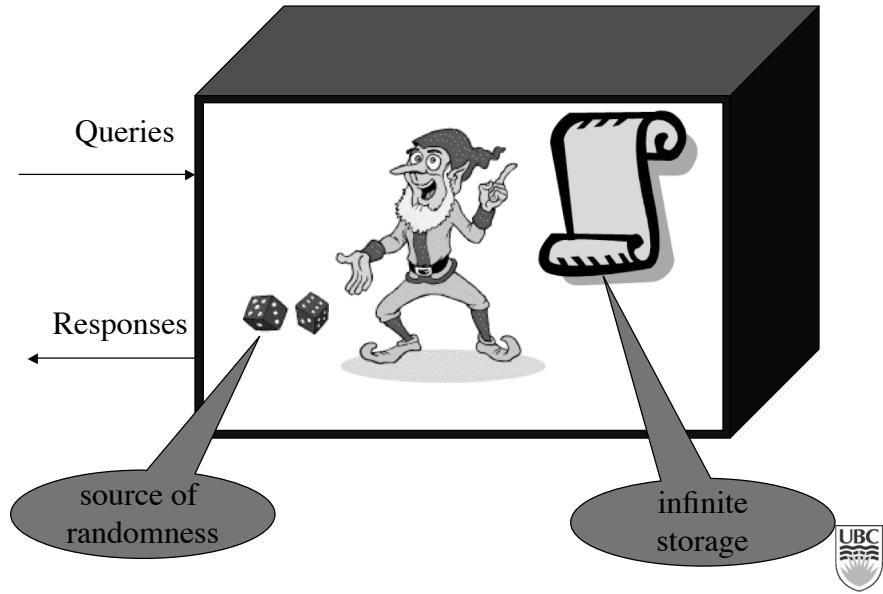


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Random Oracle Model

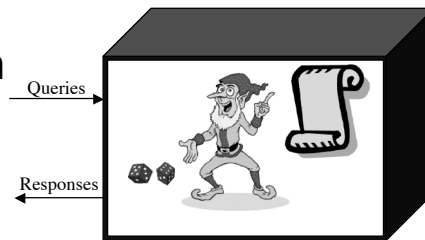
Read Anderson 5.3

What is Random Oracle Model?



Random Function as Random Oracle

- In: string of any length



- Out: random string of fixed length

- Applications:

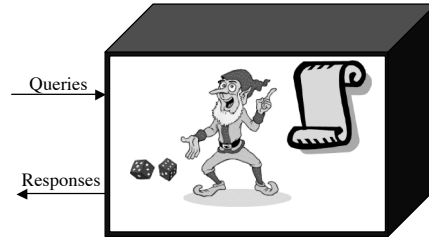
- One-way functions
- Hash functions
 - Message digests
 - Time stamping

Properties

- “One-wayness”
- No input inference from output $h(M|K)$
- Few collisions

Random Generator (Stream Cipher) as Random Oracle

- In:
 - short string (key)
 - length of the output



- Out: long random stream of bits (keystream)
- Applications:
 - Communications encryption
 - Storage encryption

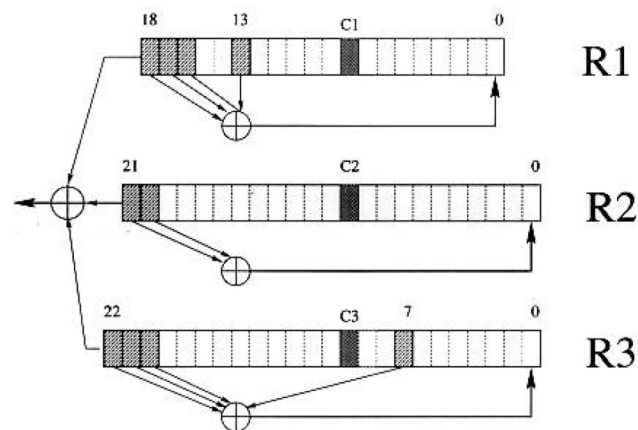
Properties

- Should not reuse
 - Use *seed*



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Example: A5 stream cipher for GSM



$$m = \text{Majority} (c_1, c_2, c_3)$$

Figure 1: The A5/1 stream cipher.

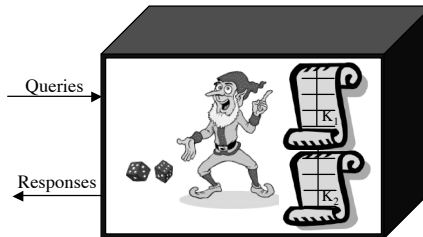


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From: Alex Biryukov, Adi Shamir, David Wagner "Real Time Cryptanalysis of A5/1 on a PC"

Random Permutation (Block Cipher) as Random Oracle

- In
 - fixed size short string (plaintext) M ,
 - DES -- 64 bits
 - Key K



- Out
 - same fixed size short string (ciphertext) C

Notation

- $C = \{ M \}_K$
- $M = \{ C \}_K$

Properties

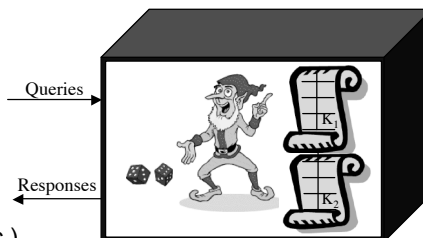
- Invertible



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Attacks on Block Ciphers

- Attack types
 - Known plaintext attack
 - Chosen plaintext attack
 - Chosen ciphertext attack
 - Chosen plaintext/ciphertext attack
 - Related key attack ($K + 1$, $K + 2$, etc.)
- Attack objectives
 - forgery attacks-- deduce the answer to the query which the attacker has not made yet
 - key recover attacks -- recover the key
- Why attack types are important?
 - DES
 - 2^{47} chosen plain texts
 - 2^{43} known plain texts

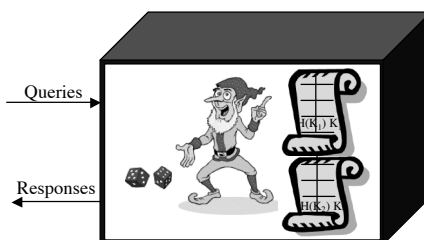


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Public Key Encryption and Trap-door One-Way Permutation as Random Oracle

- Public Key Encryption Scheme:

- Key pair (KR, KR^{-1}) generation function from random string R
 - $KR \rightarrow KR^{-1}$ is infeasible
- $C = \{M\}_{KR}$
- $M = \{C\}_{KR^{-1}}$



- In:

- fixed size short string (plaintext) M ,
- Key KR

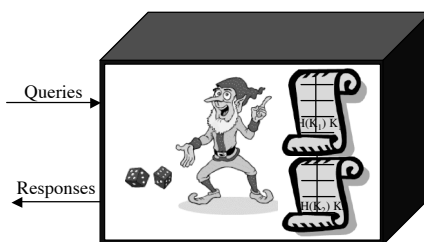
- Out: fixed size short string (ciphertext) C



Digital Signature as Random Oracle

- Public Key Signature Scheme:

- Key pair $(\sigma R, VR)$ generation function
 - $VR \rightarrow \sigma R$ is infeasible
- $S = \text{Sig}_{\sigma R}(M)$
- $\{\text{True}, \text{False}\} = \text{Ver}_{VR}(S)$



	Signing	Verifying
Input	Any string $M + \sigma R$	$S + VR$
Output	$S = \text{hash}(M) \mid$ cipher block	"True" or "False"



Summary

- Historical background
 - Caesar and Vigenère ciphers
 - One-time pad
 - One-way functions
 - Asymmetric cryptosystems
- The Random Oracle model
 - Random functions: Hash functions
 - Random generators: stream ciphers
 - Random Permutations: block ciphers
 - Public key encryption and trapdoor one-way permutations
 - Digital signatures

