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#### **Developing Secure Software**

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#### **Security Flaws**

Malicious logic

- Malware
  - Stuxnet worm
- Protection and Detection Techniques
  - Limitation of malware detection
  - Possible solutions

Non-malicious program errors

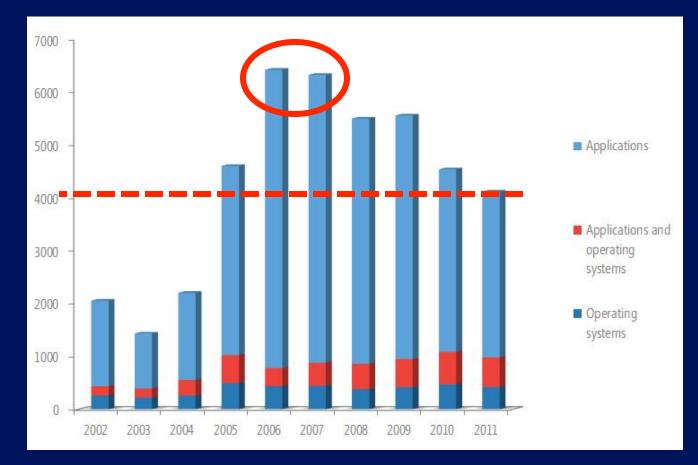
Today lecture

Last lecture



#### **Vulnerability Report Statistics**

#### Application and operating system vulnerability disclosures



(MS Security Intelligence Report, http://www.microsoft.com/ security/sir/story/default.aspx#!10year\_vulnerabilities)



#### Outline

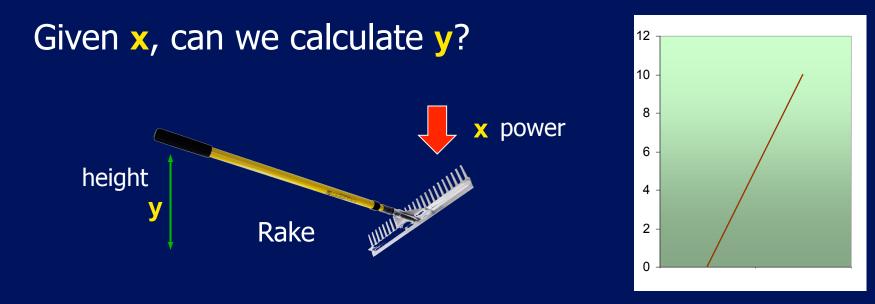
- Why developing secure software is hard?
- How are security bugs different?
- How does buffer overflow work?
- Guidelines for developing secure software





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## Why are there so many vulnerabilities in software?



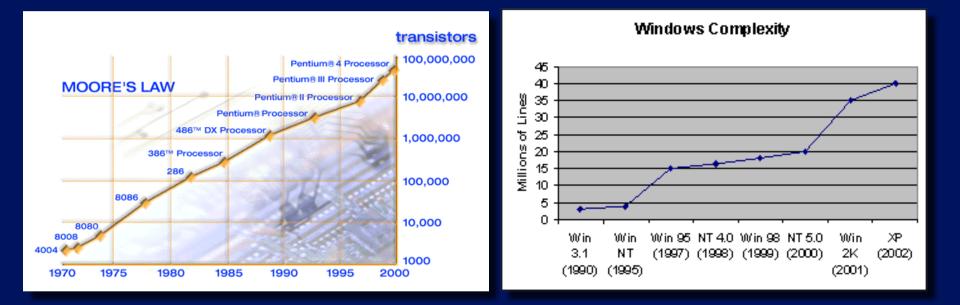
What makes this simple mechanical system predictable?

- Linearity (or, piecewise linearity)
- Continuity (or, piecewise continuity)
- Small, low-dimensional state spaces

Systems with these properties are (1) easier to analyze, and (2) easier to test.



### **Increasing complexity of computers**



Computers enable highly complex systems
 Software is taking advantages of this
 Highly non-linear behaviors; large, high-dim. state spaces



# Other software properties make security more difficult

#### **The Trinity of Trouble**

#### Connectivity

• The Internet is everywhere and most software is on it

#### Complexity

 A lot of functions: Networked, distributed, mobile, feature-full

#### Extensibility

 Systems can evolve in unexpected ways and be changed on the fly



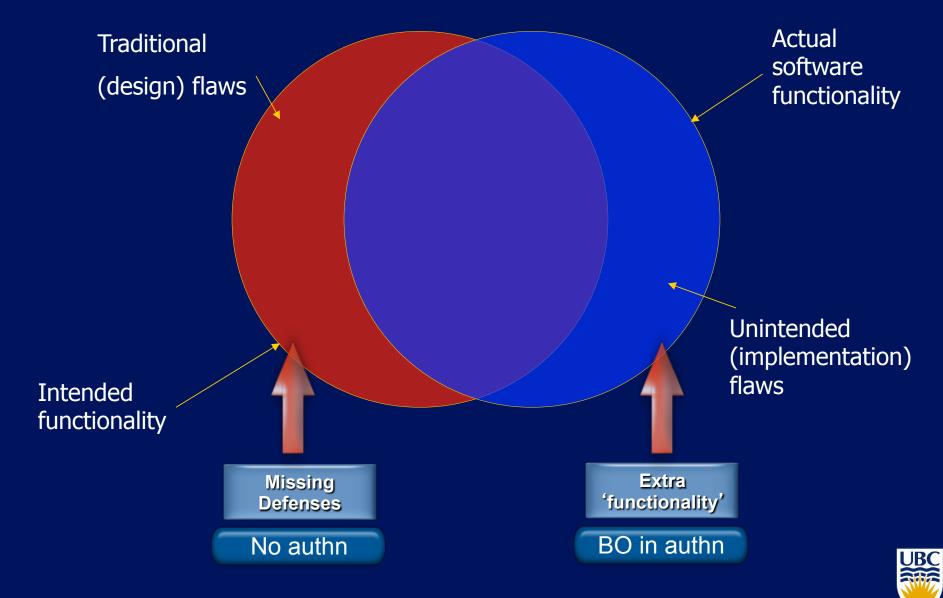


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#### How Are Security Bugs Different?

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#### Intended vs. Implemented Behavior



#### Design flaws

Incorrect

#### Weak Authentication

Supposed to do A but did B instead

Missing
 Supposed to do A *and* B but did only A.



## Security problems are complicated

#### **Implementation Flaws**

- Buffer overflow
  - String format
- Race conditions
  - TOCTOU (time of check to time of use)
- Unsafe environment variables
- Unsafe system calls
  - System()
- Untrusted input problems

#### Design Flaws

- Misuse of cryptography
- Compartmentalization problems in design
- Privileged block protection failure (DoPrivilege())
- Catastrophic security failure (fragility)
- Type safety confusion error
- Insecure auditing
- Broken or illogical access control
- Method overriding problems (subclass issues)

Which one is more frequent?

#### **Buffer Overflow**

- Can be done on the stack or on the heap.
- Can be used to overwrite the return address (to redirect the control flow of a program).



## How Buffer Overflow Works

#### Adopted from the material by Dave Hollinger

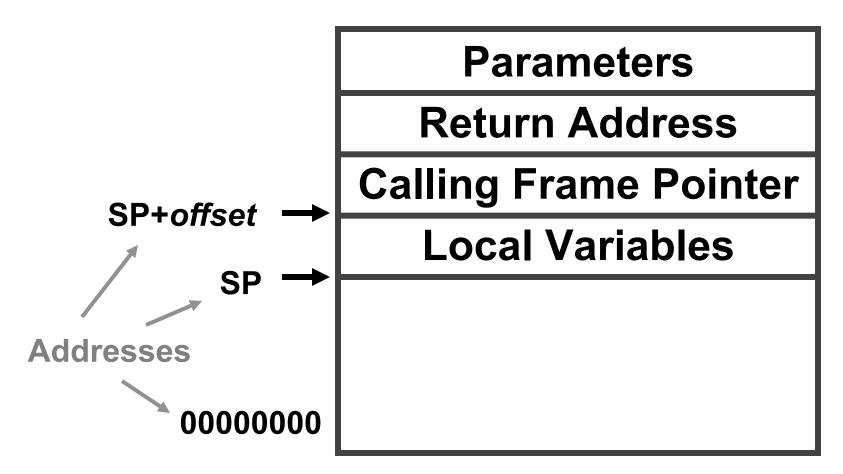
## **Overview of Buffer Overflow**

- The general idea is to overflow a buffer so that it overwrites the return address.
- When the function is done it will jump to whatever address is on the stack.
- We <u>put some code in the buffer</u> and <u>set</u>
   <u>the return address to point to it!</u>

## The Problem

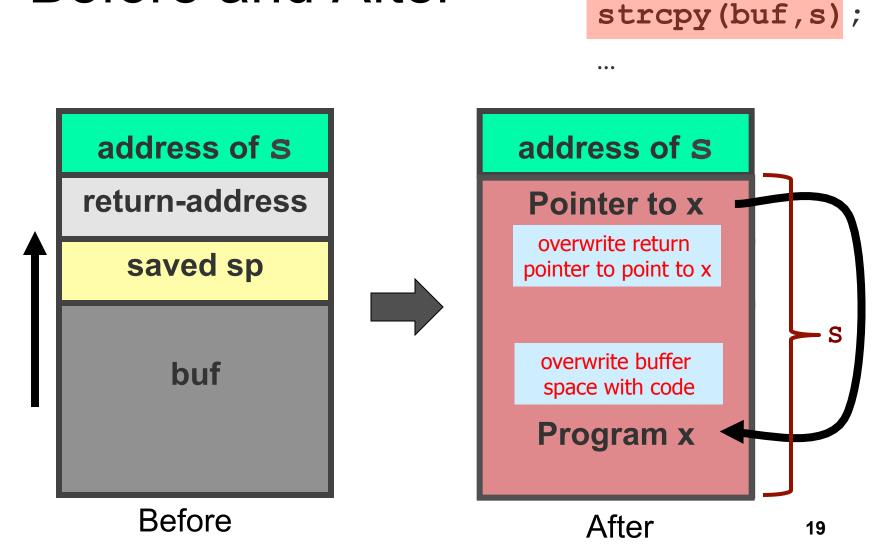
```
void foo(char *s) {
 char buf[10];
 strcpy(buf,s);
 printf("buf is %s\n",s);
foo("thisstringistolongforfoo");
```

## A Stack Frame



(SP: Stack Pointer)

Sample Stack	18 addressof(y=3) <i>return address</i> saved stack pointer y x buf
<pre>x=2; foo(18); y=3;</pre>	<pre>void foo(int j) {     int x,y;     char buf[100];     x=j;  }</pre>



void foo(char \*s) {

char buf[100];

## **Before and After**

## **Two Issues**

- 1. How do we know <u>what value the</u> <u>pointer should have</u> (the new "return address").
  - It's the address of the buffer, but how do we know what address this is?
- 2. How do we put the "program x" into the string "s"?

## An Example of Program x

#include <stdio.h>

```
char *args[] = {"/bin/ls", NULL};
```

```
void main(void) {
    execv("/bin/ls",args);
    printf("I'm not printed\n");
}
```

## Generating a String

 You can take code like the previous slide, and generate machine language.

Copy down the individual byte values and build a string.

## A Sample Program/String

Does an execv() of /bin/ls:

```
unsigned char cde[] =
\xeb\x1f\x5e
    \x89\x76\x08\x31\xc0\x88\x46\x07\x89\x4
    6\x0c\xb0\x0b\x89\xf3\x8d\x4e\x08\x8d
    \x56\x0c\xcd\x80\x31\xdb
    \x89\xd8\x40\xcd\x80\xe8\xdc\xff\xff
    \xff/bin/1s
```

We use this string for buffer overflow!

## Sample Overflow Program

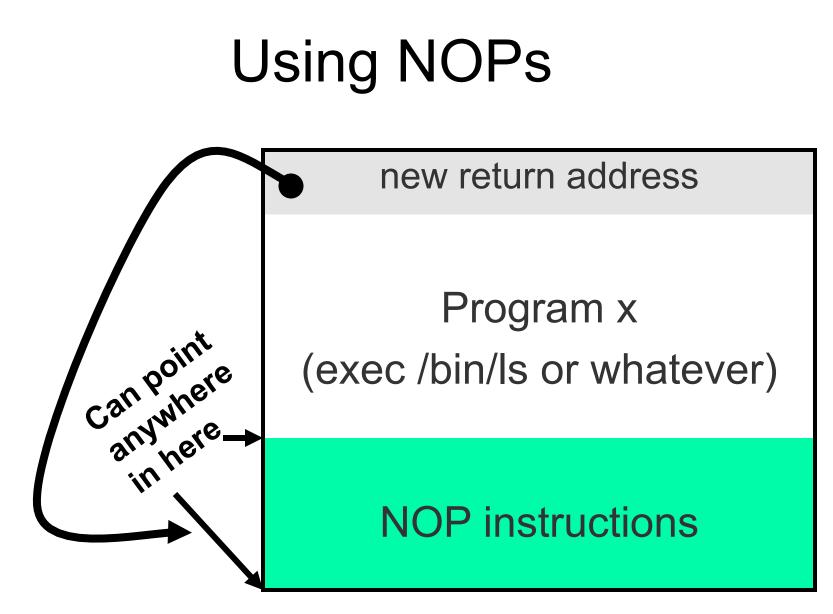
unsigned char cde[] = "\xeb\x1f\...

```
void foo(char *s) {
   char buf[10];
   strcpy(buf,s);
   printf("buf is %s\n",s);
}
```

```
int main(void) {
    printf("Running foo\n");
    foo(cde);
    printf("foo returned\n");
}
```

## **Guessing New Address**

• We need to know the address of program x.



As long as the new return-address points to a NOP we are OK

## Estimating the stack size

- We can also guess at the location of the return address relative to the overflowed buffer.
- Put in a bunch of new return addresses!

## Estimating the Location

new return address Real program

#### nop instructions

## Demo - Spock

- <u>http://nsfsecurity.pr.erau.edu/bom/Spock.html</u>
- How can you gain access as Dr. Bones when you don't know the correct password?

#### **Techniques for Preventing Buffer Overflow Attacks**

#### Write or Execute, but not both

- No program segment loaded into memory is both writable and executable
- Address Space Layout Randomization (ASLR)
  - Prevents an attacker from predicting information needed for correctly changing information flow towards the desirable computation



#### Pervasive C problems lead to bugs

#### Calls to watch out for

Instead of:	Use:
gets(buf)	fgets(buf, size, stdin)
strcpy(dst, src)	strncpy(dst, src, n)
strcat(dst, src)	strncat(dst, src, n)
<pre>sprintf(buf, fmt, a1,)</pre>	<pre>snprintf(buf, fmt, a1, n1,) (where available)</pre>
*scanf()	Your own parsing

- Hundreds of such calls
- Use static analysis to find these problems
  - ITS4, SourceScope
- Careful code review is necessary





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#### How to Develop Secure Software?

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#### Guidelines

- 1. Produce quality software
- 2. Build security into development process
- 3. Practice principles of designing secure systems
- 4. Know how systems can be compromised
- 5. Develop and use guidelines and checklists
- 6. Choose safer languages, VMs, OSs, etc.
- 7. Provide tool support

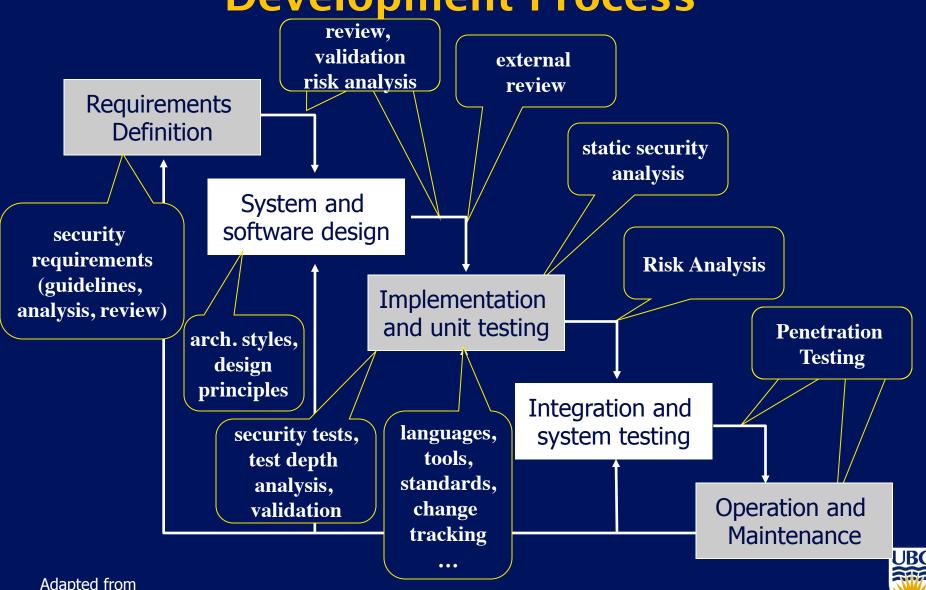


#### **1. Produce Quality Software**

Use well structured effective processes
e.g., Capability Maturity Model (CMM), \*-CMM
Use precise requirements and specifications



#### 2. Build Security into Development Process



35 D. Verdon and G. McGraw, "Risk analysis in software design," *IEEE Security & Privacy*, vol. 2, no. 4, 2004, pp. 79-84.

#### **Follow Best Practices**

- These best practices should be applied throughout the lifecycle
- Tendency is to "start at the end" (penetration testing) and declare victory
  - Not cost effective
  - Hard to fix problems
- Start as early as possible

- Abuse cases
- Security requirements analysis
- Architectural risk analysis
- Risk analysis at design
- External review
- Test planning based on risks
- Security testing (malicious tests)
- Code review with static analysis tools



# **3. Practice principles of designing secure systems**

#### **Principles of Designing Secure Systems**

- 1. Least Privilege
- 2. Fail-Safe Defaults
- 3. Economy of Mechanism
- 4. Complete Mediation
- 5. Open Design
- 6. Separation of Privilege
- 7. Least Common Mechanism
- 8. Psychological Acceptability
- 9. Defense in depth
- 10. Question assumptions



#### 4. Know How Systems Can Be Compromised

- 1. Make the Client Invisible
- 2. Target Programs That Write to Privileged OS Resources
- **3.** Use a User-Supplied Configuration File to Run Commands That Elevate Privilege
- 4. Make Use of Configuration File Search Paths
- 5. Direct Access to Executable Files
- 6. Embedding Scripts within Scripts
- 7. Leverage Executable Code in Nonexecutable Files
- 8. Argument Injection
- 9. Command Delimiters
- **10.** Multiple Parsers and Double Escapes
- **11.** User-Supplied Variable Passed to File System Calls
- **12.** Postfix NULL Terminator
- 13. Postfix, Null Terminate, and Backslash
- 14. Relative Path Traversal
- **15.** Client-Controlled Environment Variables
- **16.** User-Supplied Global Variables (DEBUG=1, PHP Globals, and So Forth)
- 17. Session ID, Resource ID, and Blind Trust
- **18.** Analog In-Band Switching Signals (aka "Blue Boxing")
- **19.** Attack Pattern Fragment: Manipulating Terminal Devices
- **20.** Simple Script Injection
- 21. Embedding Script in Nonscript Elements
- 22. XSS in HTTP Headers
- 3823. HTTP Query Strings

- 24. User-Controlled Filename
- 25. Passing Local Filenames to Functions That Expect a URL
- 26. Meta-characters in E-mail Header
- 27. File System Function Injection, Content Based
- 28. Client-side Injection, Buffer Overflow
- 29. Cause Web Server Misclassification
- **30.** Alternate Encoding the Leading Ghost Characters
- **31.** Using Slashes in Alternate Encoding
- 32. Using Escaped Slashes in Alternate Encoding
- **33.** Unicode Encoding
- 34. UTF-8 Encoding
- 35. URL Encoding
- **36.** Alternative IP Addresses
- 37. Slashes and URL Encoding Combined
- 38. Web Logs
- **39.** Overflow Binary Resource File
- **40.** Overflow Variables and Tags
- **41.** Overflow Symbolic Links
- 42. MIME Conversion
- **43.** HTTP Cookies
- 44. Filter Failure through Buffer Overflow
- **45.** Buffer Overflow with Environment Variables
- **46**. Buffer Overflow in an API Call
- **47.** Buffer Overflow in Local Command-Line Utilities
- 48. Parameter Expansion
- 49. String Format Overflow in syslog()





### **5. Develop Guidelines and Checklists**

#### Example from Open Web Application Security Project (www.owasp.org):

- Validate Input and Output
- Fail Securely (Closed)
- Keep it Simple
- Use and Reuse Trusted Components
- Defense in Depth
- Security By Obscurity Won't Work
- Least Privilege: provide only the privileges absolutely required
- Compartmentalization (Separation of Privileges)
- No homegrown encryption algorithms
- Encryption of all communication must be possible
- No transmission of passwords in plain text
- Secure default configuration
- Secure delivery
- No back doors



#### 6. Choose Safer Languages, VMs, OSs, etc.

- C or C++?
- Java or C++?
- Managed C++ or vanilla C++?
- .NET CLR or JVM?
- Windows XP or Windows 2003?
- Linux/MacOS/Solaris or Windows?



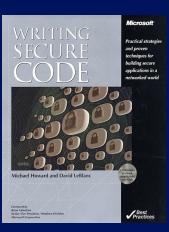
## 7. Use Good Tools

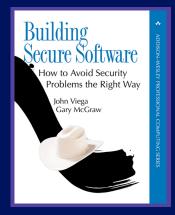
- automated tools for formal methods
  - http://www.comlab.ox.ac.uk/archive/formal-methods.html
- code analysis tools
  - RATS <u>http://www.securesw.com/rats</u>
  - Flawfinder <u>http://www.dwheeler.com/flawfinder</u>
  - ITS4 <u>http://www.cigital.com/its4</u>
  - ESC/Java <u>http://www.niii.kun.nl/ita/sos/projects/escframe.html</u>
  - PREfast, PREfix, SLAM <u>www.research.microsoft.com</u>
  - Fluid <u>http://www.fluid.cmu.edu</u>
  - JACKPOT research.sun.com/projects/jackpot
  - Many more ...

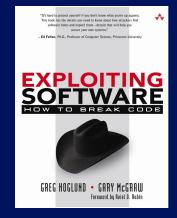


#### **Relevant Books**









#### and many more ...



#### module summary

- developing secure software is hard because it's
  - nonlinear, large, extensible, complex, has sideeffects, networked
- security bugs are different because they are undocumented side-effects
- buffer overflow works through overriding return address and replacing data with code
  guidelines for developing secure software





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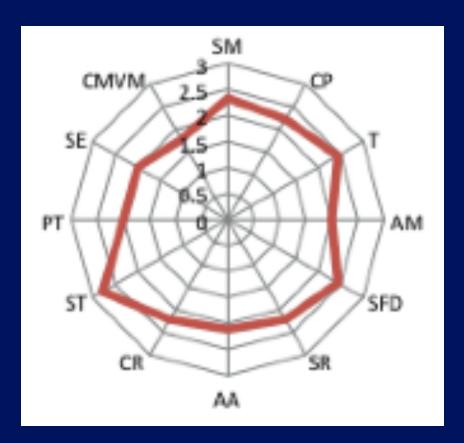
#### Case Study: Build Security In Maturity Model

#### **BSIMM**

Framework derived from SAMM Beta
Based on collected data from 9 large firms

Governance	Intelligence	SSDL Touchpoints	Deployment
Strategy and Metrics	Attack Models	Architecture Analysis	Penetration Testing
Compliance and Policy	Security Features and Design	Code Review	Software Environment
Training	Standards and Requirements	Security Testing	Configuration Management and Vulnerability Manage- ment





Source: "Building Security In Maturity Model" by Gary McGraw, Brian Chess, Sammy Migues

