

Adversary Models

CPEN 442 “Introduction to Computer Security”

Konstantin Beznosov

why we need adversary models?

attacks and countermeasures are meaningless without

elements of an adversary model

objectives

- obtain secret(s): decrypt cipher-text, guess/find password
- obtain access to assets: access to an account, full or partial control of a system or its parts

initial capabilities

- knowledge of (1) keys, passwords, and other secrets, (2) system/ environment design/architecture
- access to the system's source code and other implementation details
- partial access to a system (PC, server, mobile device)
- partial control of a system (direct browser to a URL, control of a low-privilege account)

capabilities during the attack

- passive: eavesdropping messages
- active: modifying, re-playing, or removing messages
- running code on the target system
- observing system at run-time

Dolev-Yao model

the network is completely under the adversarial control
can record, delete, replay, reroute, reorder, and completely control the scheduling of messages.

the adversary is the network

the honest participants send their messages only to the adversary and receive messages only from the adversary.

the adversary can choose the recipient and auxiliary information for its messages with total non-determinism

initial knowledge of the adversary

- the public keys (K_{Pub}),
- the private keys of subverted participants ($K_{Adv} \subseteq K_{Priv}$),
- the identifiers of the principals (I), and
- the nonces the adversary itself generates ($R_{Adv} \subseteq R$), which are assumed to be distinct from all nonces generated by honest participants.

Dolev-Yao model (continued)

message M is derivable by adversary from a set of messages S , if it's possible to produce by applying the following operations a finite number of times:

- decryption with known or learned private keys
- encryption with public keys
- pairing of two known elements
- separation of a pair into its components

Chip & PIN



EMV protocol

Europay, MasterCard, VISA (EMV) -- protocol for payment cards with chips (and PINs)

750M cards currently deployed

a three phase protocol:

Card authentication

type of card, issuer, verification method list etc)

Cardholder verification, based on verification method list,

PIN

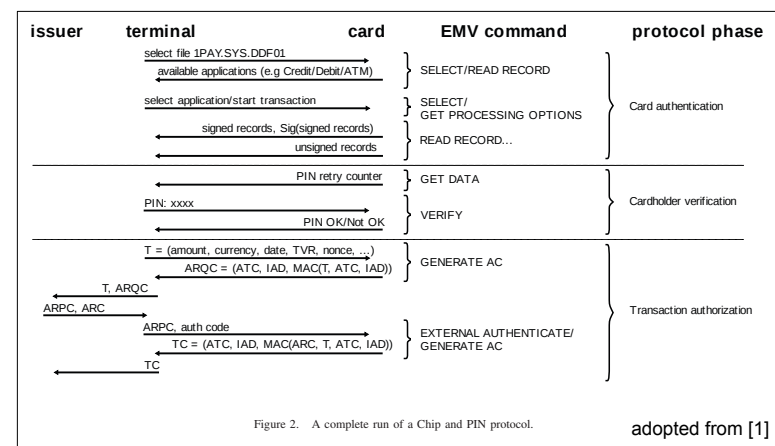
signature

nothing

Transaction authorization

card generates secured transaction info for the issuing bank clearance

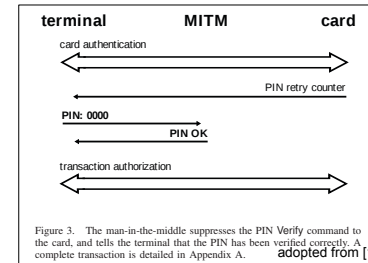
a complete run of a Chip & PIN protocol



video clip

<http://www.youtube.com/watch?v=JPAX32lgkrw>

cardholder verification step

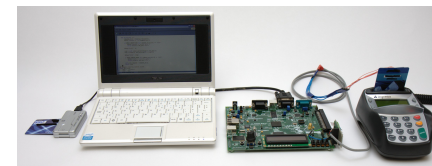
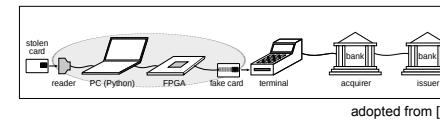


attacker tricks the card into “thinking” it’s doing a chip-and-signature transaction while the terminal “thinks” it’s chip-and-PIN.



adopted from [1]

the attack



adopted from [1]

adversary model

objectives

pay to a street merchant with a stolen payment C&P card

initial capabilities

can still payment C&P cards

can purchase or make necessary equipment for the MITM attack

capabilities during the attack

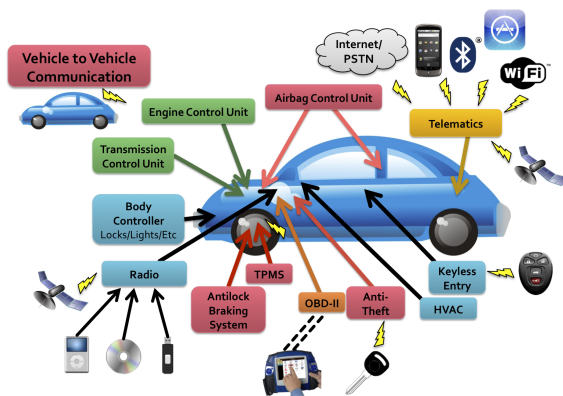
conceal the equipment from the merchant's staff

conceal the fact that the fake card has wires attached to it

insert the fake card in the merchant's terminal

Security Analysis of a Modern Car

today cars



adopted from [2]

indirect physical access: media player attack

attack 1: vestigial radio reflash from CD code

attack 2: WMA parsing bug -> buffer overflow
on-radio debugger

insert CD containing malicious WMA file
compromise the car

short-range wireless: Bluetooth attack

common embedded Bluetooth stack on telematics unit

strcpy() bug

Android trojan compromises telematics ECU

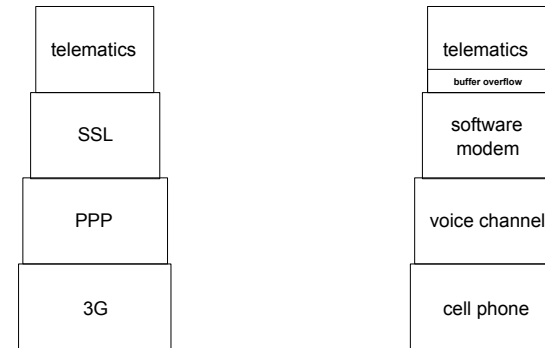
can undetectably pair a bluetooth device

USRP-based software radio

brute force PIN

cannot be unpaired with standard interface

long-range wireless: cellular attack



call telematics unit

transmit malicious payload (using modem protocol or just play malicious sound track over phone)

what's next?

remotely trigger code from prior compromise

proximity trigger

broadcast trigger (FM RDS)

short-range targeted trigger (Bluetooth)

global targeted trigger (cellular)

what can an adversary do with this?

car theft

compromise car

locate it via GPS

unlock doors

start engine

bypass anti-theft

video demo: <http://www.youtube.com/watch?v=bHfOzilwXic> (minute #16)

surveillance

compromise car

continuously report GPS coordinates

stream audio recorded from the in-cabin mic

adversary model

objectives

take control over parts or the whole car in order to perform surveillance, theft, or cause car accident.

initial capabilities

access to equipment and documentation to develop and test an attack
extract device's firmware
reverse engineer firmware
identify and test vulnerable code paths
weaponize exploits

capabilities during the attack (one of the three)

indirect physical access to the car
interacts with a physical object that interacts with the car
diagnostic tool that plugs directly into OBD-II port
entertainment systems (CD player, digital multimedia port, iPod Out)
short-range wireless signals (between 5 and 300 meters)
Bluetooth, Remote Key Entry, RFID car keys, Tire Pressure Monitoring Systems, WiFi, Dedicated Short Range Communications
long-range wireless signals (greater than 1 km)
broadcast channels: GPS, satellite radio, digital radio, Radio Data System, Traffic Message Channel

summary: adversary model

objectives

initial capabilities

capabilities during the attack

references

Chip and PIN is Broken, Murdoch, Steven J.; Drimer, Saar; Anderson, Ross; Bond, Mike; , "Chip and PIN is Broken," 2010 IEEE Symposium on Security and Privacy (SP), pp.433-446, 16-19 May 2010, doi: 10.1109/SP.2010.33

"Comprehensive Experimental Analyses of Automotive Attack Surfaces," S. Checkoway, D. McCoy, B. Kantor, D. Anderson, H. Shacham, S. Savage, K. Koscher, A. Czeskis, F. Roesner, T. Kohno, USENIX Security, August 10–12, 2011.